

**ENM 202**  
**OPERATIONS RESEARCH (I)**  
**OR (I)**  
**5**

**LECTURE NOTES**  
**SPECIAL CASES in the SIMPLEX**  
**METHOD APPLICATIONS**

# SPECIAL CASES in the SIMPLEX METHOD APPLICATIONS

1)DEGENERACY	2)ALTERNATIVE OPTIMA	3)UNBOUNDED SOLUTION	4)INFEASIBLE SOLUTION
<p>1)An LP is degenerate if it has at least one bfs in which a BV is equal to zero</p> <p>2) A degenerate optimal solution occurs when the model has at least one redundant constraint</p> <p>3) Degeneracy can occur temporarily or permanently.  <u>When degeneracy occurs permanently</u>, the objective function value will never improve and the procedure will repeat a sequence of iterations.we call this as cycling or circling.  <u>When degeneracy occurs temporarily</u>, the simplex iterations must always be continued until the last iteration satisfies the optimality condition</p>	<p>1) When the objective function is parallel to a binding constraint, alternative optimal solution case occurs.</p> <p>2) In the simplex tableau of the final iteration, if the coefficient of a NBV in z-row is equal to zero then this indicates that alternative optimal solution exists.</p>	<p>1) In some LPs the values of variables may be increased indefinitely without violating any of the constraints, meaning that the solution space is unbounded.</p> <p>2) As a result, the objective value may increase (max. case) or decrease (min. case) indefinitely. In this case both the solution space and the optimal objective value are unbounded.</p> <p>The general rule for recognizing unboundedness</p> <p>a) If at any iteration, all the constraint coefficients of any NBV are zero or negative, then the solution space is unbounded in that direction.</p> <p>b) In addition, If the objective coefficient of that variable is negative in max problem, or positive in min problem then the objective value is also unbounded.</p>	<p>1) LP with inconsistent constraints has no feasible solutions</p> <p>2) This situation never occur if all the constraints are of the type <math>\leq</math> (assuming nonnegative RHSs) since the slack variables always provide a solution.</p> <p>3)For other type of constraints (<math>=, \geq</math>) we use artificial variables. If the model has one artificial variable with a positive value in the optimum iteration then the model has no solution and infeasible solution case occur</p>

# 1)DEGENERACY

- 1)An LP is degenerate if it has at least one bfs in which a BV is equal to zero
- 2) A degenerate optimal solution occurs when the model has at least one redundant constraint
- But there are no reliable techniques for identifying the redundant constraints directly from the simplex tableau.

- Example
- (Degenerate Optimal Solution- degeneracy occurs temporarily))

- |                             |  |
|-----------------------------|--|
| • LP: Max $z = 3x_1 + 9x_2$ | LP in Standard Form: Max $z = 3x_1 + 9x_2$ |
| • St                        | st   |
| • $x_1 + 4x_2 \leq 8$       | $x_1 + 4x_2 + x_3 = 8$                     |
| • $x_1 + 2x_2 \leq 4$       | $x_1 + 2x_2 + x_4 = 4$                     |
| • $x_1, x_2 \geq 0$         | $x_1, x_2, x_3, x_4 \geq 0$                |

# Degenerate Optimal Solution

Iteration	Basic	x1	x2	x3	x4	RHS	MRT
(0)	z	-3	-9	0	0	0	
x2 enters x3 leaves	x3	1	4	1	0	8	$(8/4)=2^*$
	x4	1	2	0	1	4	$(4/2)=2^*$
(1)	z	-3/4	0	9/4	0	18	
	x2	1/4	1	1/4	0	2	$2/(1/4)=8$
x1 enters x4 leaves	x4	1/2	0	-1/2	1	0	$0/(1/2)=0^*$
(2)	z	0	0	3/2	3/2	18	
optimum	x2	0	1	1/2	-1/2	2	
	x1	1	0	-1	2	0	

# Grafical Representation of Degenerate Optimal Solution

$$\text{Max } z = 3x_1 + 9x_2$$

s.t.

$$x_1 + 4x_2 \leq 8$$

$$x_1 + 2x_2 \leq 4$$

$$x_1, x_2 \geq 0$$

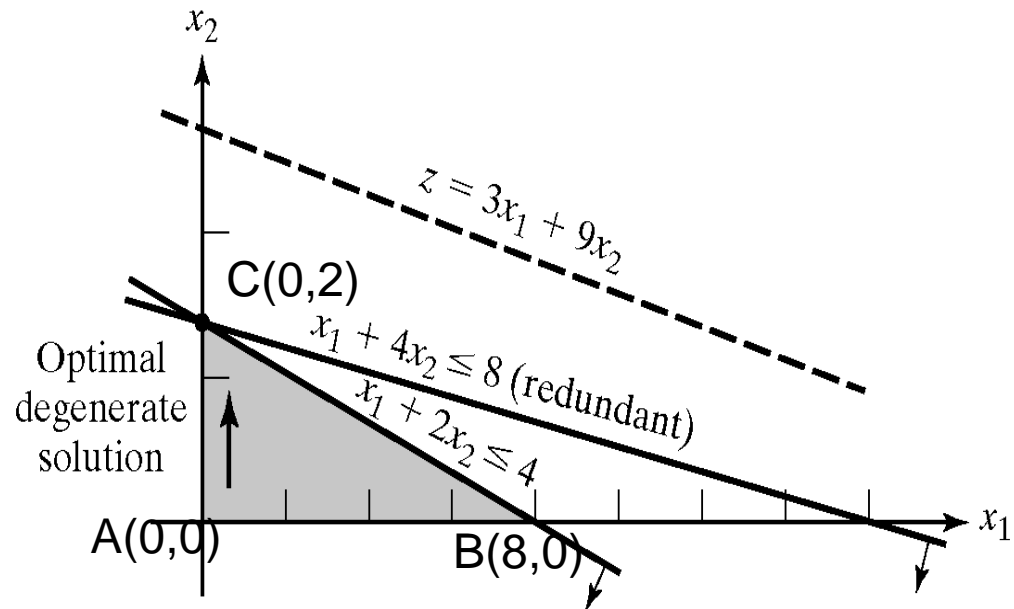


Figure 3.7

LP degeneracy in Example 3.5-1.

# CYCLING (CIRCLING)

- Example
- (Taha OR Book 7th ed., Set 3.5A, Problem 4, pp 106,)
- LP:
- $\text{Max } z = (3/4)x_1 - 20x_2 + (1/2)x_3 - 6x_4$
- St
- $(1/4)x_1 - 8x_2 - x_3 + 9x_4 \leq 0$
- $(1/2)x_1 - 12x_2 - (1/3)x_3 + 3x_4 \leq 0$
- $x_3 \leq 1$
- $x_1, x_2, x_3, x_4 \geq 0$
- You will notice that the starting all-slack bfs at iteration 1 will appear identically in iteration 7.

## 2)ALTERNATIVE OPTIMA INFINITY of SOLUTIONS

- Example
  - LP: Max  $z = 2x_1 + 4x_2$
  - St
  - $x_1 + 2x_2 \leq 5$
  - $x_1 + x_2 \leq 4$
  - $x_1, x_2 \geq 0$
- LP in Standard Form: Max  $z = 2x_1 + 4x_2$   
st  
 $x_1 + 2x_2 + x_3 = 5$   
 $x_1 + x_2 + x_4 = 4$   
 $x_1, x_2, x_3, x_4 \geq 0$

## 2)ALTERNATIVE OPTIMA INFINITY of SOLUTIONS

Iteration	Basic	x1	x2	x3	x4	RHS	MRT
(0)	z	-2	-4	0	0	0	
x2 enters x3 leaves	x3	1	2	1	0	5	5/2*
	x4	1	1	0	1	4	4/1
(1) optimum	z	0	0	2	0	10	
	x2	1/2	1	1/2	0	5/2	5
x1 enters x4 leaves	x4	1/2	0	-1/2	1	3/2	3*
(1') alt. optimum	z	0	0	2	0	10	
	x2	0	1	1	-1	1	
	x1	1	0	-1	2	3	



## 2. Alternative Optimal Solutions (Infinity of Solutions) (Multiple Optimal Solutions)

Example (Taha OR Book pp 106)

LP: Max  $Z = 2x_1 + 4x_2$

St

$$x_1 + x_2 \leq 4 \quad (1)$$

$$x_1 + 2x_2 \leq 5 \quad (2)$$

$$x_1 \geq 0, x_2 \geq 0 \quad (3,4)$$

B(0,5/2)

C(3,1)

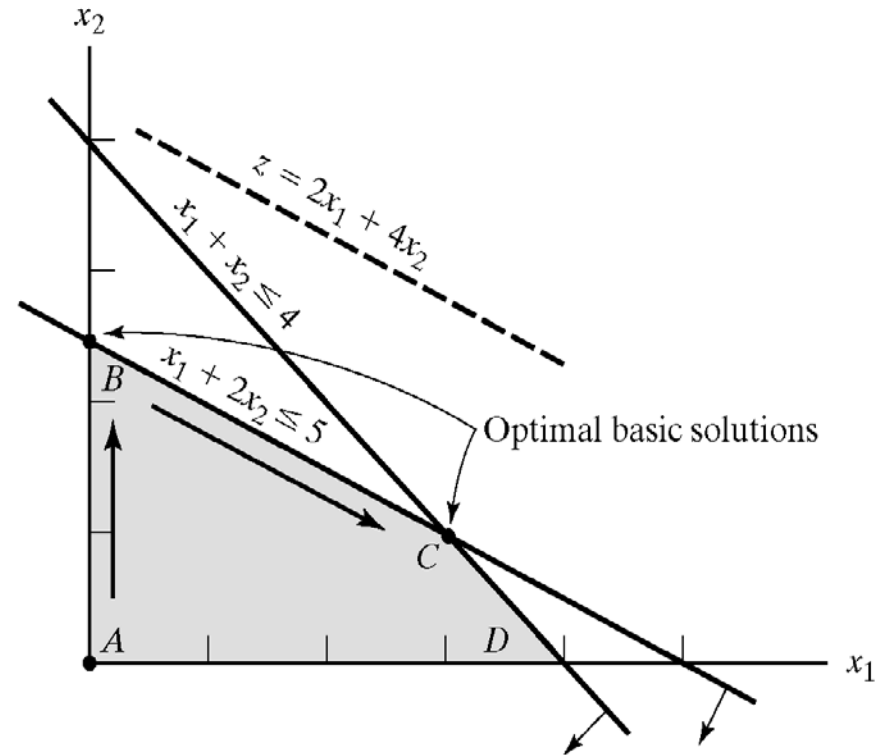


Figure 3.9

LP alternative optima in Example 3.5-2.

### 3) UNBOUNDED SOLUTION

- Example (Unbounded Objective Value)

- Example

- LP: Max  $z = 2x_1 + x_2$

- St

- $x_1 - x_2 \leq 10$

- $2x_1 \leq 40$

- $x_1, x_2 \geq 0$

LP in Standard Form: Max  $z = 2x_1 + x_2$

st

$x_1 - x_2 + x_3 = 10$

$2x_1 + x_4 = 40$

$x_1, x_2, x_3, x_4 \geq 0$

# (Unbounded Objective Value)

Iteration	Basic	x1	x2	x3	x4	RHS	MRT
(0)	z	-2	-1	0	0	0	
x1 enters x3 leaves	x3	1	-1	1	0	10	10/1*
	x4	2	0	0	1	40	40/2

- In the starting Tableau both x1 and x2 are candidates for entering the solution.
- Because x1 has the most negative coefficient, it is selected as the entering variable.
- However, all the constraint coefficients under x2 are negative or zero, meaning that x2 can be increased indefinitely without violating any of the constraints.

Each unit increase in x2 will increase z by 1, an infinite increase in x2 will also result in an infinite increase in z. Thus the problem has no bounded solution.

# (Unbounded Objective Value)

Iteration	Basic	x1	x2	x3	x4	RHS	MRT
(0)	z	-2	-1	0	0	0	
x1 enters x3 leaves	x3	1	-1	1	0	10	10/1*
	x4	2	0	0	1	40	40/2
(1)	z	0	-3	2	0	20	
	x1	1	-1	1	0	10	ignore
x2 enters x4 leaves	x4	0	2	-2	1	20	20/2*
(2)	z	0	0	-1	3/2	50	
	x1	1	0	0	1/2	20	
	x2	0	1	-1	1/2	10	

# Graphical Representation of Unbounded Solution Case

•Example (Taha OR Book pp 109)

$$\text{Max } Z = 2x_1 + x_2$$

St

$$x_1 - x_2 \leq 10 \quad (1)$$

$$2x_1 \leq 40 \quad (2)$$

$$x_1, x_2 \geq 0 \quad (3,4)$$

IN A CORRECTLY  
MODELLED LP, IT  
SHOULD NOT OCCUR

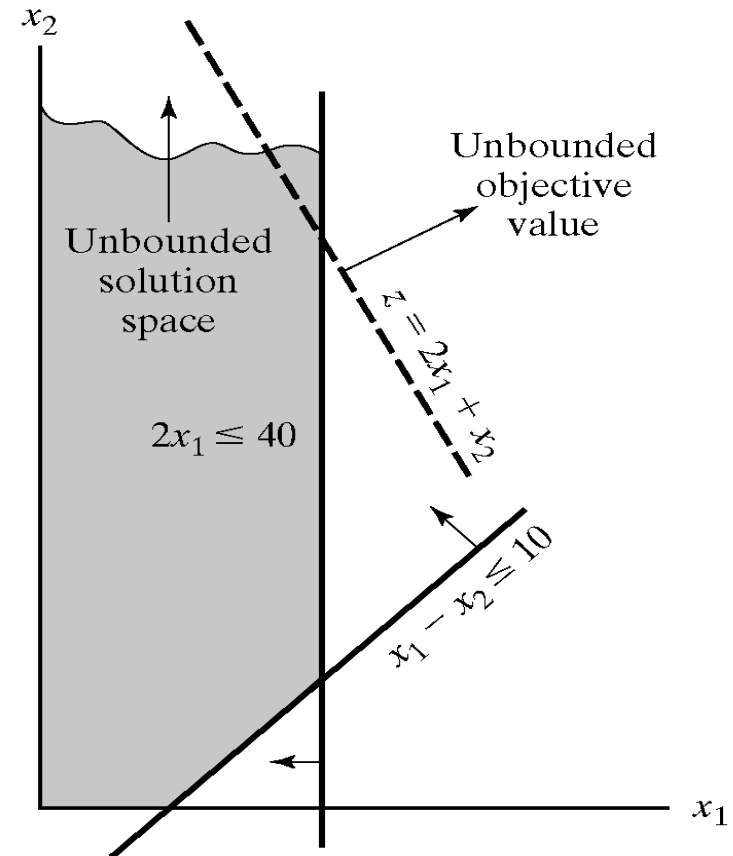


Figure 3.10

LP unbounded solution in Example 3.5-3.

# UNBOUNDED SOLUTION SPACE

- Example (Unbounded Solution Space but Finite Optimum Objective Value)

- LP: Max  $z = 6x_1 - 2x_2$

- St

- $2x_1 - x_2 \leq 2$

- $x_1 \leq 4$

- $x_1, x_2 \geq 0$

LP in Standard Form: Max  $z = 6x_1 - 2x_2$

st

$2x_1 - x_2 + x_3 = 2$

$x_1 + x_4 = 40$

$x_1, x_2, x_3, x_4 \geq 0$

# (Unbounded Solution Space but Finite Optimum Objective Value)

Iteration	Basic	x1	x2	x3	x4	RHS	MRT
(0)	z	-6	2	0	0	0	
x1 enters x3 leaves	x3	2	-1	1	0	2	$2/2=1^*$
	x4	1	0	0	1	4	$4/1=4$
(1)	z	0	-1	3	0	6	
	x1	1	-1/2	1/2	0	1	ignore
x1 enters x3 leaves	x4	0	1/2	-1/2	1	3	$3/(1/2)=6^*$
(2)	z	0	0	2	2	12	
optimum	x1	1	0	0	1	4	
	x2	0	1	-1	2	6	

## 4) INFEASIBLE SOLUTION

- LP models with inconsistent constraints have no feasible solutions
- Example (Taha OR Book pp 111)
- LP: Max  $Z = 3x_1 + 2x_2$   
St  
 $2x_1 + x_2 \leq 2$
- $3x_1 + 4x_2 \geq 12$
- $x_1 \geq 0, x_2 \geq 0$
- LP in Standard Form with  
Artificials:
- Max  $Z = 3x_1 + 2x_2$   
St  
 $2x_1 + x_2 + x_3 = 2$
- $3x_1 + 4x_2 - x_4 + x_5 = 12$   
 $x_1, x_2, x_3, x_4, x_5 \geq 0$



# Example for INFEASIBLE SOLUTION

Iteration	Basic	x1	x2	x4	x3	x5	RHS	MRT
Inconsistent tableau	z	-3	-2	0	0	M	0	
	x3	2	1	0	1	0	2	
	x5	3	4	-1	0	1	12	
(0)	z	-3-3M	-2-4M	M	0	0	-12M	
x2 enters x3 leaves	x3	2	1	0	1	0	2	2/1*
	x5	3	4	-1	0	1	12	12/4
(1)	z	1+5M	0	M	2+4M	0	4-4M	
PSEUDO OPTIMUM	x2	2	1	0	1	0	2	
	x5	-5	0	-1	-4	1	4	

# Graphical Representation of Infeasible solution case

Example (Taha OR Book pp 111)

$$\text{Max } Z = 3x_1 + 2x_2$$

St

$$3x_1 + 4x_2 \geq 12 \quad (1)$$

$$2x_1 + x_2 \leq 2 \quad (2)$$

$$x_1 \geq 0, x_2 \geq 0 \quad (3,4)$$

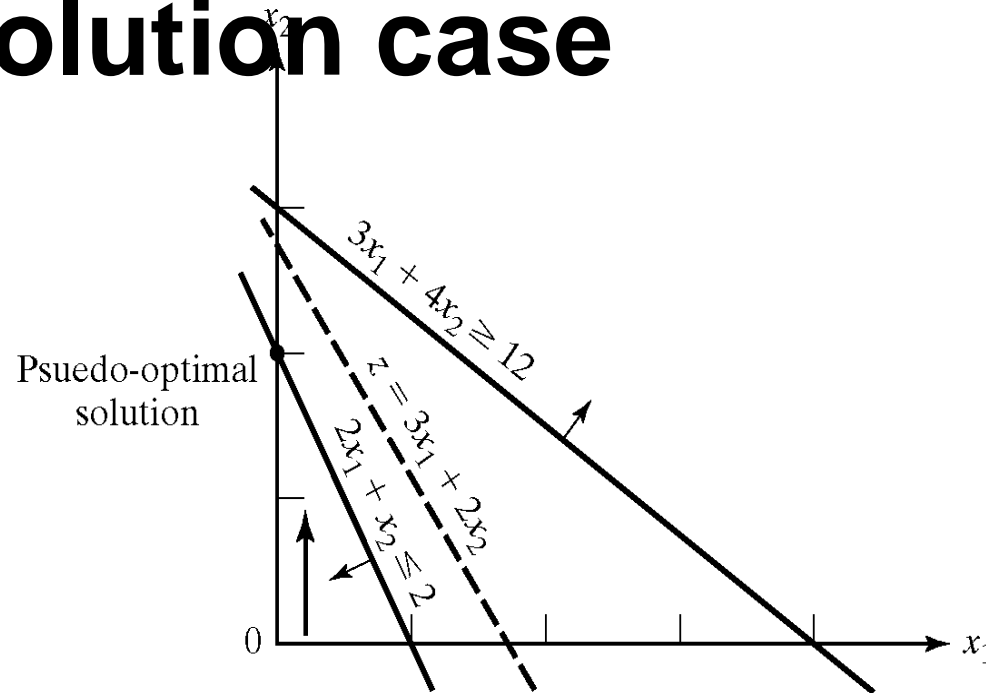


Figure 3.11

Infeasible solution of Example 3.5-4.